
Hua Wang and Zhijun Zheng
School of Information and Electronic Engineering, Zhejiang University of Science and Technology, Hangzhou, China
Email: wanghua96@gmail.com, zjzheng9999@sina.com

Abstract—There had been past attempts at making adaptation strategies with analytic model that used the current environment information in Cyber-physical System (CPS) to keep the software architecture from deteriorating. However, the research is still in its infancy on the closely-related issue of how to take corrective self-adaptive actions to reconcile the CPS system behavior with the variability of the survival environment. In particular, architects have almost no assistance in reasoning about questions such as: How should we stage the architectural evolution to improve self-adaptation to accommodate uncertainties in CPS? In this work, the specification of software architecture is extended using CHAM (Chemical Abstract Machine) in the presence of uncertainty. The key benefits of our approach are that it leverages standard software architecture models, and quantifies behaviors within the system in terms of relevant architectural elements. Compared to the traditional model, the proposed method could arrange optimized response sequence and adjust the behaviors within the software system to be adaptive to the new situations over time in CPS.

Index Terms—self-adaptive, software architecture, cyber-physical system, uncertainty, chemical abstract machine

I. INTRODUCTION

Cyber–Physical System (CPS) is the next generation of engineered systems in which computing, communication, and control technologies are tightly integrated. CPSs are integrations of computation and physical processes [1]. To satisfy emerging and ever-changing user requirements and expectations, software systems in CPS should be self-adaptive to preserve the system architecture from deterioration. Endowing such a system with an adaptive property can strengthen the survival ability of the system, which we term collective self-adaptation in the presence of cyber-physical convergence. Recent retrofitting autonomous adaptation capabilities into software-intensive systems are prevalently done in an ad-hoc fashion. The proposed approaches should offer a compositional and systematic model using software architecture, which guides the topology of the constituent computational elements (such as components and connectors) of the system. The behavior of the whole system is the synergistic effect of these interacting elements. Problems occur when a component is used in an unanticipated way and/or components become a constituent of a larger complex system. The dynamic context where the system is running can cause unexpected behavior resulting from a combination of the unanticipated use of components and interactions between components, which stems from the fact that CPSs are inherently dynamical and uncertain. Dealing with context uncertainty and predicting complex behavior in social networking systems are challenged [2].

To this end, a recent common approach to monitor and adapt system behavior at runtime is to decouple self-adaptive mechanisms from the target system by considering the intrinsic nature of CPS. The non-invasive manners have the main advantage of that style-specific adaptation knowledge has not been “hardwired” into the target system, i.e., realizing separation of concerns. However, a variety of sources of uncertainty are introduced while utilizing these separate control units [3]. These uncertainties may include bringing the system into unknown system states, the random user requests for services, the unanticipated effects of performing adaptation policies, to name a few. There is a substantial body of discussion literatures on handling uncertainty of software behaviors, such as [4-6], etc.

There have also been past attempts at making adaptation strategies with analytic model that used current environment information, such as [7-9]. However, on the closely-related issue of how to take corrective self-adaptive actions to reconcile the CPS system behavior with the variability of the survival environment, the research is still in its infancy. An adaptation agent was employed to collect massive distributed data from widespread sensors, and then constructed and executed adaptation plans [10]. Evolutionary computation was applied to mitigate uncertainty in the case of high-assurance adaptive software, especially in CPS [11]. Jeffrey and colleagues put forward an architecture evolution means among potential candidate evolution paths [12].

© 2014 ACADEMY PUBLISHER
doi:10.4304/jcp.9.4.802-811
The increasing shift toward unanticipated adaptation calls for revisiting current reactive self-adaptive approaches to anticipate situations on the recent future and adjust the behaviors within the system under consideration in CPS. The traditional method puts itself to some extent that the adaptation process can only be preplanned and defined in a limited space. The need to anticipate the aforementioned adaptation tasks has driven the development of predictive self-adaptation. The unanticipated inherence and complexity of upcoming services and applications make proactive self-adaptation essential. In particular, architects have almost no assistance in reasoning about questions such as: How should we stage the architectural evolution to improve self-adaptation to accommodate uncertainties in CPS? This presents the main technical challenge, i.e., the proposed model can learn from the environment of cyber-physical convergence. In this paper, we propose a collective self-adaptive method based on the software architecture framework in the presence of uncertainty attributes in CPS. Compared to the traditional model, the proposed method could arrange optimized response sequence and adjust the behaviors within the software system to be adaptive to the new situations over time.

In our previous work [13], Hidden Markov Model was employed to model uncertainty and learn from history behavior of target system, and then anticipatory actions were issued on the target system. In this paper, Chemical Abstract Machine (CHAM) [14] is employed as a tool to specify software architecture. CHAM is good as the description of the dynamics and parallels of systems. The semantic of the dynamic operation of software architecture is described and analyzed by formal methods of CHAM. Beneficial extension of dynamic features of software architecture is presented by inspecting uncertainty. Formal reasoning of specification of self-adaptive software architecture is realized. The focus of the software architecture specification by CHAM is how to describe the dynamic interaction behavior of components comprising of software architecture. In this work, the proactive ability is added by underlying anticipatory self-adaptation. Our approach is novel as it leverages standard software architecture models, and quantifies behaviors within the system in terms of relevant architectural elements. The new contribution of this paper is to apply self-adaptive software architecture in CPS field base on our previous paper [15]. We extended the model check of CHAM and specified Safety and Liveness properties of CHAM by performing two different algorithms. Also, the new experiment was designed to show the efficiency of the proposed method and more results were obtained.

The remainder of this paper is organized as follows: Section II introduces the specification of collective self-adaptive software architecture based on CHAM. We discuss the uncertainty in our context in Section III and model check of the specification of collective self-adaptive software architecture in Section IV. In Section V, the improved Viterbi-I algorithm is presented. In Section VI, we report on a case study on the self-adaptive method based on software architecture in CPS followed by discussion of the remaining challenges and possible directions in future research.

II. SPECIFICATION OF COLLECTIVE SELF-ADAPTIVE SOFTWARE ARCHITECTURE

From the formal semantics of CHAM, a conclusion can be reached that the reaction rule is an ideal tool to reflect the dynamic behaviors within the target system. Components of software architecture could be considered as molecules to represent the inner status of CHAM. Relevant molecules can interact with each other by different statuses. The specification of architectural CHAM includes architecture (molecules) grammar of static components, the grammar of the initial state of software architecture and the syntax of the system in response to the dynamic development phase according to the grammar rules. The initial status set is the subset of relevant molecules configured by grammar. The initial static configuration of software architecture relates to the initial status set of CHAM. The transform rules applied to initialize the solution set define how the system develops dynamically. The initial solution is relevant to the static configuration of the targeted system. Each molecule of the initial solution corresponds to the initial status of each architectural element. The transform rules define how software architecture evolves dynamically from the initial configuration.

A. Specification of Collective Software Architecture based on CHAM

Definition 1. The specification of collective self-adaptive software architecture based on CHAM is defined as:

\[ Cham = \left( ComM, ConnM, R, S, s_0, S_f, Path \right) \]

where \( ComM \) is the grammar set of component molecules of software architecture; \( ConnM \) is the grammar set of connector molecules of software architecture; \( R \in WP \) is the reaction rules of a molecule set; \( S \) is the solution set; \( s_0 \in S \) is the initial solution set of software architecture; \( S_f \in S \) is the terminal solution set of software architecture; and \( Path \) is the path of status change of software architecture.

The component molecule set reflects component ontology while connector molecule set represents the role set of connectors. Components, connectors and constraints are basic elements of software architecture. Port and Role are other basic elements of software architecture. The interface is the only bridge for the interaction of components with other elements of software architecture. The port of components is the communication point for components and software survival environment. The interface is composed of a group of ports. Connectors have interfaces that are composed of a group of roles. The role represents the actor participating the activity. In this sense, ports, interfaces and roles imply interaction and constraint of architectural elements, which belongs to grammar layer. On the other hand, \( Path \) demonstrates how connectors use
Weaving Policy (WP) to instruct the dynamic behaviors from the initial solution to terminal solution, and this is belongs to semantic layer. WP is built by extensible ECA rules based on our previous works [16].

Since Cham considers the initial status of the target system as the initial solution, the system is a single flow launched by an initial solution [17]. However, molecules, membranes and sub solutions concurrently react each other probably, as a result, a special connection symbol “∥” is used to handle the concurrent reaction. The elements connected by “∥” can react concurrently, but there is no intersection. Consequently, the molecule algebra $\Sigma_{Cham}$ is redefined as follows:

**Definition 2.** $\Sigma_{Cham}$ is defined as:

$$Molecule := Com | Connector | Connector \odot M | M \odot Connector | M \odot M | M | M | Cmd | Connection$$

$Com := ID : Type$

$Type := AC | PC$

$Adv := Role | DUP (number) : Molecule | ID : Mediator \odot WP$

$Role := RoleName : Molecule | RoleName = RoleName = ID$

$Cmd := be (Com) | ec (ID) | re (Com) | ac (Connector) | ec (ID) |

$Connector := RoleName - RoleName$

where $Com$ is components identified by $ID$; $AC$ is Aspect Component; $PC$ is Primary Component; $Connector$ is composed of $Mediator$ and $WP$ (Mediator reasons about $WP$ by the symbol ‘⊗’); the command of $be$ (build a component), $ec$ (erase a component), $uc$ (unify components) and $sc$ (split a component) can be used to realize the static evolution of software architecture. Component encapsulating crosscutting concern (such as security, requirement and distribution) is defined as $AC$ while component performing a functional operation is defined as $PC$. $Connector$ coordinates the $AC$ and $PC$.

**Definition 3.** Cluster component molecule is represented by component molecule and is encapsulated by a membrane, defined as follows:

$$ClusterCom := \{Com_0, Com_1, \ldots, Com_l\}$$

**Definition 4.** The initial solution of Cham is the initial static configuration of software architecture and the starting point of the evolution of software architecture. Suppose that the initial sub solution is $s_0 := s_0_1 \cdots s_0_n$, then the initial solution of software architecture is:

$$s_0 = s_0_1 || \ldots || s_0_n, s_1 = \left\{m_0, m_1, \ldots, m_l\right\}$$

**B. Reaction Rules of Cham**

**Definition 5.** Reaction rules are defined as follows:

$$T_1 = m_1 \rightarrow m_1 \rightarrow m_1$$

$$T_2 = \beta (ac : AC), pc : PC, pc' : PC \rightarrow ac : AC, pc : PC, pc' : PC$$

$$T_3 = \beta (PC, ac : AC), pc : PC, pc' : PC \rightarrow pc : PC, ac : AC, pc' : PC$$

$$T_4 = ec (pc : PC), pc : PC = ec (pc : PC)$$

$$T_5 = ec (ac : AC), pc : PC = ec (ac : AC)$$

$$T_6 = \rho (ac : AC), pc : PC \rightarrow ac : PC, pc : PC$$

$$T_7 = \rho (ac : AC), pc : PC \rightarrow ac : PC, pc : PC$$

$$T_8 = \rho (ac : AC), pc : PC \rightarrow ac : PC, pc : PC$$

where operator ‘⊗’ represents the connection based on roles for components and connectors. Architectural elements can be modeled by a group of PCs, ACs, roles and WPs. If a certain solution does not permit the reaction of specified molecules, membrane and airlock reaction rules are used.

**Definition 6.** Membrane reaction rules:

$$T_{10} = \{ac \odot pc_1 \left\langle \{pc_2, \ldots, pc_k\} \right\rangle \rightarrow ac \odot pc_1 \left\langle \{pc_2, \ldots, pc_k\} \right\rangle\}$$

where $ac \in AC, pc_1, \ldots, pc_k \in PC$, and $\{pc_2, \ldots, pc_k\}$ represents membrane and its reaction rules are similar to solution. Molecules can be picked up from solution by ‘<’.’ If a reaction does not reflect the inner status of membrane, the following reaction rule is required.

**Definition 7.** Airlock reaction rule:

$$T_{11} = \{ac \odot pc_1 \left\langle \{pc_2, \ldots, pc_k\} \right\rangle \rightarrow ac \odot pc_1 \left\langle \{pc_2, \ldots, pc_k\} \right\rangle\}$$

where $ac \in AC, pc_1, \ldots, pc_k \in PC$.

**III. Uncertainty in Cham**

### A. Source of Uncertainty

As mentioned before, randomness inherent in the dynamic context and the noisy nature of the predictive model leads to uncertainty in CPS. The model must explicitly dispose the uncertainty. Some research results constitute a contribution to deal with the uncertainty, such as Poladian [18] proposed an approach to self-adaptation that leverages predictions of future resource availability. However, unlike dealing with uncertainty in resource predictions, we present uncertainty as probability distributions, i.e., consider the uncertainty of user requests for services as a stochastic process. The arrival of requests into a service and the transition from Request to Request is uncertain, and may be represented as a Poisson distribution and transition probability matrix with some parameters, respectively.
B. Uncertainty Analysis

Uncertainty is presented by probability distribution to analyze and analog software architecture. For example, the arrival-rate of requests is considered as a Poisson distribution with certain mean and variance. The accurate analysis of distribution helps to analyze the dynamic behavior within the system under consideration meaningfully. A particular operator AssDis (Assign Distribution) is used to capture a probability distribution. The operator can be used as a command in-built. Software architect uses the operator to specify the uncertainty of software survival environment.

Definition 8. Uncertainty Cham:

\[
\text{Cham} := \text{Cnd} | \text{AssDis}
\]

\[
\text{AssDis} := \text{pc} : \text{PC} \otimes (\text{dt} : \text{DistributionType})
\]

\[
\text{DistributionType} := \left\{ \text{Normal, Exponential, Poisson,...} \right\} \bullet \left[ \frac{\text{prop}_1 = \text{val}_1,}{\text{prop}_2 = \text{val}_2, \ldots} \right]
\]

where \(\otimes\) represents the type of probability distribution of components in software survival environment, such as normal distribution, exponential distribution and Poisson distribution. The distribution attributes can be expressed by \(\bullet\). For example, an online health evaluation system makes an assessment of the diabetes risk of a patient. The assessment component is DiabetesPC. Suppose the requests of assessment of the diabetes risk satisfy Poisson distribution, the probability feature can be added into Cham specification, expressed by AssDis as follows:

\[
\text{AssDis} := \text{DiabetesPC} \otimes (\text{Poisson} \bullet \lambda = 0.23)
\]

By using the extensive probability distribution, Cham can employ some standard probability technique to evaluate the effect of uncertainty on software architecture and then self-adaptive actions can be performed. Uncertainty lies in the stochastic nature within the context. Cham focuses on request of users to think about self-adaptive actions. Self-adaptive model considers uncertainty as a probability distribution, that is to say, requests of users are regarded as a stochastic process. The arrival of requests is uncertain and also the transmission of requests (from Request to Request) is also uncertain. Accordingly, history request information and relevant probability distribution are speculated. The transition probability matrix can be used to represent the transition of requests.

Our proposed method provides with a user-defined probability property for designers to specify the new probability type that captures different distribution characteristics in different domains. A Poisson distribution, for example, can be specified to depict the arrival rate of user requests as illustrated in Figure 1. According to the definition of the Poisson distribution type in Cham, model designer can use the self-definition type to specify certain desired probability characteristics in different domains.

### Figure 1. Specifying poisson distribution type in Cham.

Then, uncertainty is inserted into the specification of software architecture. Software architect can describe statistical features of system behavior employing probability distribution to specify architectural behavior in a wide range. All consequent research can be done based on the mathematical analysis model. The uncertain evolution process of software architecture can be depicted by this way. After that, the future features of the target system can be conferred. The self-adaptation is a specific nature of the behavior of software architecture. The driving of self-adaptation of software architecture is the survival environment in CPS. On the other hand, the self-adaptation of the software-intensive system is closely relative to software architecture. Consequently, the uncertainty of software survival environment and self-adaptability of software architecture are the two base point of our work. By analyzing uncertainty of software survival environment, we can design self-adaptive software architecture to induce the evolution within the system and realize the self-adaptation to survival environment.

IV. MODEL CHECK OF CHAM

A. Export Specification of Software Architecture

The advantage of Cham model is to effectively describe dynamic evolution of software architecture. The dynamic features of Cham model can be described by LTS (Labeled Transition System) [19]. Cham has no the dedicated mechanism to control reaction rules at every time because it is possible for Cham to apply several rules at the same time. A certain rule is selected nondeterministically among these rules. The formalism of Cham allows to verify and validate multiple properties by employing different analysis tools and approaches. Especially, LTS is deduced from Cham using its operationabilities and further reasoning is feasible in LTS. Altruistic lock [20] is used to export the status tree of LTS from Cham model.

Considering the long transaction for the reaction of molecules resulting in postponing the reaction of other molecules, altruistic lock allows molecules of long transaction to release relevant locks once these molecules need no more data to permit other molecules to react. The export algorithm uses generalized list to describe the data structure for reaction rules and solution. The altruistic lock protocol ensures the parallel reaction to generate the right results. The protocol maintains two transition rule sets \(L(c)\) and \(D(c)\), where \(c\) is a certain molecule, defined as follows.

Definition 9. \(L(c)\) means transition rule sets performing a Lock operation for \(c\). \(D(c)\) means transition rule sets performing Donate operation.
And the protocol maintains another two transition rule sets, \( W(l) \) and \( P(l) \) defined as follows.

**Definition 10.** \( W(l) \) denotes transition rule sets of which donate molecule sets include the total access sets of \( l \). \( P(l) \) denotes molecules required by \( l \) are locked.

When solutions are initial, there is \( P(l) = L(c) = \emptyset \). The Lock, Unlock and Denote algorithm is illustrated in Fig. 2.

\[ \text{Figure 2. Lock, Unlock and Denote algorithms.} \]

In Lock algorithm, \( pl \) denotes the transition rule sets where \( l \) only follows up and \( wd \) denotes the transition rule sets which follows up. Thus, the transition rules with alturistic lock can entirely follow up other transition rules. The export algorithm for Cham model is designed illustrated in Figure 3.

\[ \text{Algorithm 4. ChamPS4 generates LTS} \]

\[
\begin{align*}
1 & \text{TRACE=NUll; \slash/initialize} \\
2 & s=\text{initial solution}; \slash/initial solution \\
3 & s_f=\text{final solution}; \slash/terminal solution \\
4 & \text{while}(\text{true}) \\
5 & \text{isSol}() \text{while}(\text{true}) \\
6 & \text{terminate, } \text{if the initial solution is equal with terminal one, end} \\
7 & \text{break} \\
8 & \text{false} \\
9 & \text{Matc}(c_i, c_j, T_i, T_j) \text{while}(\text{true}) \\
10 & \text{if}(Z_i = \text{NUll}) \\
11 & \text{else} \\
12 & \text{TRACEEnd}(c_i, c_j, T_i, T_j) \slash/\text{generate new solution} \\
13 & \text{construct reaction tree, } V \slash/\text{construct reaction tree} \\
14 & \text{end} \\
15 & \text{end} \\
16 & \text{end} \\
18 & \text{if}(\text{TRACE} = \text{null}) \text{continue} \\
19 & \text{end}
\end{align*}
\]

\[ \text{Figure 3. Export algorithm for Cham model.} \]

The first line of the algorithm \textbf{TRACE} is used to follow the track of every pair of \( c_i, c_j \) and the reaction rule \( T_i \). The ninth line denotes that molecule \( c_i, c_j \) and reaction rule \( T_i \) are matched by \textbf{Match}(c_i, c_j, T_i). The 19th line judges if molecule \( c_i \) and reaction rule \( T_i \) are locked. And furthermore, if molecule \( c_i \) is not locked, molecule \( c_i, c_j \) and reaction rule \( T_i \) are added into \textbf{TRACE} and repeats it.

**B. Specify Safety and Liveness Properties of Cham using LTS**

Modal mu-calculus has excellent logic description and can be used to formulate the safety and liveness properties of concurrent systems. As a very expressive propositional temporal logic, modal mu-calculus is a modal logic with external fixed points employed to specify safety and liveness properties of Cham represented by as the resulting LTS. The greatest fixed operator of mu-calculus denotes the safety property of Cham suggesting that no unanticipated results are brought about whether the system has the safety property, that is, the safety property excludes a group of bad characteristics. While the least fixed point operator indicates a group of liveness attributes and specifies a special status implementing good characteristics.

The model check of LTS is performed by modal mu-calculus. The key idea is to describe the behaviors of Cham using LTS and the system properties using logic formula \( mmf \) in modal mu-calculus, respectively. As a result, whether the system holds the promising characteristics or not is converted into a mathematical problem, and that is whether \( lts \) is a model of the formula \( mmf \) or not by defining \( lts \rightarrow mmf \), which is interpreted by LTS as follows.

\[
\begin{align*}
& [Z]e = e(Z) \\
& [\ell_1 \cup \ell_2]e = [\ell_1]e \cup [\ell_2]e \\
& [\ell_1 \cap \ell_2]e = [\ell_1]e \cap [\ell_2]e \\
& [\ell_k](s) = s | s \ell_k s' \Rightarrow s' \Rightarrow [\ell_k]e \\
& [k](s) = s \forall s' s \ell k s' : s' \Rightarrow [\ell_k]e
\end{align*}
\]

where \( e \) is an environment and defined as \( Var \rightarrow 2^S \). \( Var \) is a definite variable set and \( S \) the state set of LTSs, and \( k \in K \) is an action specified in the self-adaptive language in our previous work [15]. We call \( S_j \) is the derivative of \( S_i \) on the condition \( i \lliff \)

\[
\begin{align*}
& s_i \frac{\downarrow}{\rightarrow} s_j \\
& Z_n = \ell_n
\end{align*}
\]

The mutually recursive equational block [21] is employed for the purpose of compensating for the defect that \( mmf \) can only describe the behavior of a limited part of the system attributes. An equational block has one of two forms: \( \text{min}\{E\} \) or \( \text{max}\{E\} \), where \( E \) is a list of equations as follows.

\[
\begin{align*}
& Z_i = \ell_i \\
& \vdots \\
& Z_n = \ell_n
\end{align*}
\]

where \( \ell_i \) is a basic formula, and \( Z_i \) is different from each other. Let \( Z[i] \) and \( C[i] \) be bit array and counter array, respectively, to store the information during the process of checking models. \( s.C[i] \) is true if \( s \) belongs to the set of propositional variables \( Z_i \). If \( b \) is a max block, where \( b \in B \) and \( B \) are the set of mutually recursive equational blocks, then the following possibilities exist for \( C[i] \).

1) If \( Z_i = Z_j \cup Z_k \) is an equation and belongs to \( B \), \( s.C[i] \) counts all disjunctive terms (true in terms of \( s \)) appearing on the right-hand side to the equation.  
2) If \( Z_i = \{k\} Z_j \) is in \( B \), \( s.C[i] \) counts \( l \) derivative of \( s \) associated with \( Z_j \).  
3) Others, \( C[i] \) is not used. In the same way, if \( b \) is a min block, the following possibilities exist for \( C[i] \).
1) If $Z_i = Z_j \cap Z_k$ is an equation and belongs to $B$, $s.C[i]$ counts all conjunctive items (false in terms of $s$) appearing on the right-hand side to the equation.
2) If $Z_i = (k)Z_j$ is in $B$, $s.C[i]$ counts $l$ derivative of $s$ being not associated with $Z_j$.
3) Others, $C[i]$ is not used.

The block processing about max is organized in algorithms as follows:

**Algorithm 5. Max Block Processing**

<table>
<thead>
<tr>
<th>Line</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>$B_{\text{max}} =$ initial max block; //initial max block</td>
</tr>
<tr>
<td>2</td>
<td>$\text{SVP}[m]= $ initial state variable pairs;</td>
</tr>
<tr>
<td>3</td>
<td>if $(Z_i = Z_j \cup Z_k)$ or $(Z_i = Z_\ell \cup Z_k)$ //beginning of the first case</td>
</tr>
<tr>
<td>4</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>5</td>
<td>if(Z is the left-hand side of $B_{\text{max}}$)</td>
</tr>
<tr>
<td>6</td>
<td>$s.C[j] =$;</td>
</tr>
<tr>
<td>7</td>
<td>}</td>
</tr>
<tr>
<td>8</td>
<td>if (s. C[j] = 0) //no disjunctive term meets s</td>
</tr>
<tr>
<td>9</td>
<td>delete s from $S_i$</td>
</tr>
<tr>
<td>10</td>
<td>$s.Z[j] =$false;</td>
</tr>
<tr>
<td>11</td>
<td>$\text{SVP}[m++] = s.Z[j]$;</td>
</tr>
<tr>
<td>12</td>
<td>}</td>
</tr>
<tr>
<td>13</td>
<td>}</td>
</tr>
<tr>
<td>14</td>
<td>} //end of the first case</td>
</tr>
<tr>
<td>15</td>
<td>if $(Z_i = Z_j \cup Z_k)$ or $(Z_i = Z_\ell \cup Z_k)$ //beginning of the second case</td>
</tr>
<tr>
<td>16</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>17</td>
<td>if(Z is the left-hand side of $B_{\text{max}}$) and (s. Z[j] = true)</td>
</tr>
<tr>
<td>18</td>
<td>$s.Z[j] =$false;</td>
</tr>
<tr>
<td>19</td>
<td>$\text{SVP}[m++] = &lt;s.Z[j]&gt;;</td>
</tr>
<tr>
<td>20</td>
<td>}</td>
</tr>
<tr>
<td>21</td>
<td>}</td>
</tr>
<tr>
<td>22</td>
<td>} //end of the second case</td>
</tr>
<tr>
<td>23</td>
<td>if $(Z_i = &lt;k&gt; Z_j)$ //beginning of the third case</td>
</tr>
<tr>
<td>24</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>25</td>
<td>if (Z is the left-hand side of $B_{\text{max}}$)</td>
</tr>
<tr>
<td>26</td>
<td>if (s' is s derivation on l)</td>
</tr>
<tr>
<td>27</td>
<td>for (each counter) C[j] =;</td>
</tr>
<tr>
<td>28</td>
<td>}</td>
</tr>
<tr>
<td>29</td>
<td>if (C[j] = 0) //s' has no l derivation meeting $Z_i$</td>
</tr>
<tr>
<td>30</td>
<td>s'.Z[j] =false;</td>
</tr>
<tr>
<td>31</td>
<td>$\text{SVP}[m++] = &lt;s'.Z[j]&gt;;</td>
</tr>
<tr>
<td>32</td>
<td>}</td>
</tr>
<tr>
<td>33</td>
<td>}</td>
</tr>
<tr>
<td>34</td>
<td>} //end of the third case</td>
</tr>
<tr>
<td>35</td>
<td>if $(Z_i = [k]Z_j)$ //beginning of the fourth case</td>
</tr>
<tr>
<td>36</td>
<td>for (each state s')</td>
</tr>
<tr>
<td>37</td>
<td>if (Z[i] = false) and (Z[j] is the left-hand side of $B_{\text{max}}$)</td>
</tr>
<tr>
<td>38</td>
<td>if (s' is s derivation on l)</td>
</tr>
<tr>
<td>39</td>
<td>$Z[i] =$false;</td>
</tr>
<tr>
<td>40</td>
<td>$\text{SVP}[m++] = &lt;s'.Z[j]&gt;;</td>
</tr>
<tr>
<td>41</td>
<td>}</td>
</tr>
<tr>
<td>42</td>
<td>}</td>
</tr>
<tr>
<td>43</td>
<td>} //end of the fourth case</td>
</tr>
</tbody>
</table>

**Figure 4. Algorithm for Max block processing.**

Actually, the algorithm for max block processing continuously deletes <s, Z> from the list $\text{SVP}[n]$ until $\text{SVP}[n]$ is empty as illustrated in Figure 4. At this point, the bit vector for each state relevant to the max block has the final value of fixed-point. The max block excludes the bad characteristics and expresses the safety property.

In the same way, the algorithm for min block processing is illustrated in Figure 5, which continuously deletes <s, Z> from the list $\text{SVP}[n]$ until $\text{SVP}[n]$ is empty. At this point, the bit vector for each state relevant to the min block has the final value of fixed-point. Whether a certain state meets the given formula or not can be judged by inspecting the final bit vector, that is, if <s, Z> is equal 1, then Z meets s; otherwise, if <s, Z> is equal 0, the state s has no path meeting K.

**Algorithm 6. Min Block Processing**

<table>
<thead>
<tr>
<th>Line</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>$B_{\text{min}} =$ initial max block; //initial min block</td>
</tr>
<tr>
<td>2</td>
<td>$\text{SVP}[m]= $ initial state variable pairs;</td>
</tr>
<tr>
<td>3</td>
<td>if $(Z_i = Z_j \cup Z_k)$ or $(Z_i = Z_\ell \cup Z_k)$ //beginning of the first case</td>
</tr>
<tr>
<td>4</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>5</td>
<td>if(Z is the left-hand side of $B_{\text{max}}$) and (s. Z[j] = true)</td>
</tr>
<tr>
<td>6</td>
<td>$s.Z[j] =$true;</td>
</tr>
<tr>
<td>7</td>
<td>}</td>
</tr>
<tr>
<td>8</td>
<td>if (s. C[j] = 0) //all right-hand disjunctive terms meet s</td>
</tr>
<tr>
<td>9</td>
<td>add s into $S_i$</td>
</tr>
<tr>
<td>10</td>
<td>$s.Z[j] =$true;</td>
</tr>
<tr>
<td>11</td>
<td>$\text{SVP}[m++] = &lt;s.Z[j]&gt;;</td>
</tr>
<tr>
<td>12</td>
<td>}</td>
</tr>
<tr>
<td>13</td>
<td>}</td>
</tr>
<tr>
<td>14</td>
<td>} //end of the first case</td>
</tr>
<tr>
<td>15</td>
<td>if $(Z_i = Z_j \cup Z_k)$ or $(Z_i = Z_\ell \cup Z_k)$ //beginning of the second case</td>
</tr>
<tr>
<td>16</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>17</td>
<td>if(Z is the left-hand side of $B_{\text{max}}$) and (s. Z[j] = false)</td>
</tr>
<tr>
<td>18</td>
<td>$s.Z[j] =$true;</td>
</tr>
<tr>
<td>19</td>
<td>$\text{SVP}[m++] = &lt;s.Z[j]&gt;;</td>
</tr>
<tr>
<td>20</td>
<td>}</td>
</tr>
<tr>
<td>21</td>
<td>}</td>
</tr>
<tr>
<td>22</td>
<td>} //end of the second case</td>
</tr>
<tr>
<td>23</td>
<td>if $(Z_i = &lt;k&gt; Z_j)$ //beginning of the third case</td>
</tr>
<tr>
<td>24</td>
<td>for (each $Z_i$)</td>
</tr>
<tr>
<td>25</td>
<td>if (Z is the left-hand side of $B_{\text{max}}$)</td>
</tr>
<tr>
<td>26</td>
<td>if (s' is s derivation on l)</td>
</tr>
<tr>
<td>27</td>
<td>for (each counter) C[j] =;</td>
</tr>
<tr>
<td>28</td>
<td>}</td>
</tr>
<tr>
<td>29</td>
<td>if (C[j] = 0) //s' has no l derivation meeting $Z_i$</td>
</tr>
<tr>
<td>30</td>
<td>s'.Z[j] =true;</td>
</tr>
<tr>
<td>31</td>
<td>$\text{SVP}[m++] = &lt;s'.Z[j]&gt;;</td>
</tr>
<tr>
<td>32</td>
<td>}</td>
</tr>
<tr>
<td>33</td>
<td>}</td>
</tr>
<tr>
<td>34</td>
<td>} //end of the third case</td>
</tr>
<tr>
<td>35</td>
<td>if $(Z_i = [k]Z_j)$ //beginning of the fourth case</td>
</tr>
<tr>
<td>36</td>
<td>for (each state s')</td>
</tr>
<tr>
<td>37</td>
<td>if (Z[i] = false) and (Z[j] is the left-hand side of $B_{\text{max}}$)</td>
</tr>
<tr>
<td>38</td>
<td>if (s' is s derivation on l)</td>
</tr>
<tr>
<td>39</td>
<td>$Z[i] =$true;</td>
</tr>
<tr>
<td>40</td>
<td>$\text{SVP}[m++] = &lt;s'.Z[j]&gt;;</td>
</tr>
<tr>
<td>41</td>
<td>}</td>
</tr>
<tr>
<td>42</td>
<td>}</td>
</tr>
<tr>
<td>43</td>
<td>} //end of the fourth case</td>
</tr>
</tbody>
</table>

**Figure 5. Algorithm for Min block processing.**

Actually, the max block processing results expresses the liveness property because good characteristics are included.

© 2014 ACADEMY PUBLISHER
V. HMM MODEL AND VITERBI-I ALGORITHM

In previous work [13], we address the situations of the evolution of software and focus on dynamic self-adaptation at runtime using the idea of HMM (Hidden Markov Model). We classify the modes of requested services, analyze the statistical properties of requests, and then use an XML file describing the context at runtime. The model treats every Request Type as a state and the invocation number of Components as the observation sequence. Based on this analysis, we model this environment and schedule clients’ requests to respond to clients in a more efficient and rapid fashion. However, the classic Viterbi algorithm requires much multiplication results in the loss of precision and the overflow. This will lead to a reduced performance of the algorithm. The improved Viterbi-I algorithm uses LOG function for p. The multiplication of the probability values of p can be converted to the addition of log p. The observation value probability is sorted given certain request sequence. It is magnificent for improvement of Viterbi algorithm [22] because there is no need to search full status sequence. Suppose q(i,l) denotes the optimized request path excluding l observation values.

\[
q(i,l) = \arg\max_{q_i} \prod_{t=2}^{l} \log p(q_t | q_{t-1}) + \chi_i(O_t | q_i)
\]

where \( \chi_i(O_t | q_i) = \prod_{i=j+1}^{r} \log p(o_i | q_i) \).

At time t, \( \delta(i) \) denotes the max probability generating the path sequence \( O = O_1, O_2, \ldots, O_T \) along with path \( R_1, R_2, \ldots, R_t \), i.e.,

\[
\delta(i) = \max_{\theta_i, q_1, \ldots, q_{t-1}} P(q_1, q_2, \ldots, q_t, O_t = \theta_i, O_1, O_2, \ldots, O_t / \lambda)
\]

The algorithm of the optimized request path is as follows:

1. **Initialization:**
   \( \delta(0) = \log p_i + \log b_1(q_1), 1 \leq i \leq N \)
   \( \varphi(i) = 0.1 \leq i \leq N \)

2. **Recursion:**
   for \( 2 \leq t \leq T \) and \( 1 \leq j \leq N \),
   \( \delta(j) = \max_{1 \leq i \leq N} [\delta(i,j-1) + \log a_{ij} \log b_j(O_t)] \)
   \( \varphi(j) = \arg\max_{1 \leq i \leq N} [\delta(i,j-1) + \log a_{ij}], 2 \leq t \leq T, 1 \leq j \leq N \)

3. **Termination:**
   \( p^* = \max_{1 \leq i \leq N} [\delta(i,T)] \)
   \( q^*_T = \arg\max_{1 \leq i \leq N} [\delta(i,T)] \)

4. **Solution of request sequence:**
   for \( t = T - 1, T - 2, \ldots, 1 \),
   \( q^*_t = q^*_{t+1}(q^*_t) \)

VI. CASE STUDY

To demonstrate the usefulness of our approach, we show how it can be applied to a motivating industry project to illustrate the proposed modeling method, providing real-time Pervasive Health Management Director (PHmD) system based on the medical health knowledge database by analyzing health data collected from remote wireless health devices, such as Blood Pressure Monitor, Electrocardiogram Measurement Instrument, Blood-glucose Meter and so on. A possible scenario is the realtime physiological signal is relayed to PHmD while the jogger is running with the wearable textiles or devices. PHmD accommodates for the health evaluation service based on QoS (Quality of Service) requirements subscribed by different customers who receive relevant services to monitor and improve their health level before potential diseases occur, as illustrated in Figure 6.

![Figure 6. PHmD computing environment for fitness application.](image)

PHmD will ultimately lead in the application of low-cost and innovative technological systems to daily support and assist patients, to prevent and control the ongoing of the pathology, to adjust drug therapies, and to avoid hospitalization. In particular, measurement of body movement and physical activity by inertial sensors has disclosed a wide variety of applications in favor of monitoring of bio-signals (in particular, heart rate, respiration, and gait).

An authorized healthcare professional will be able to access the health information for a customer in real-times, send health promoting messages and actuate implanted drug delivery in situ. The fitness application of PHmD computing environment uses facilities to identify disease processes or to detect adverse health-related events. The services provided by the application measure the continuous behavior of these systems, rather than their average value over some time period, to detect the dynamics.

A. Software Architecture Specification Model of PHmD

As an illustrative motivating example, let us consider the following requirement scenario for our PHmD. As a part of PHmD system requirements, Jogging Monitor specifies the Device (Wearable Textiles) communicating with the Node (Health App.) hosted in wireless cellphones, shown in Table I according to IEEE Recommended Practice for Software Requirements Specifications (SRS).

<table>
<thead>
<tr>
<th>Priority</th>
<th>Scenario Description</th>
<th>Stimulus/Response Sequences</th>
</tr>
</thead>
<tbody>
<tr>
<td>High</td>
<td>Monitor the vital sign of the jogger, the enority, if</td>
<td>1. While jogging, the wearable device is on.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>2. The data of Blood Pressure, Body Temperature, Respiratory Frequency</td>
</tr>
</tbody>
</table>

TABLE I. DESCRIPTION OF JOGGING MONITOR USING SRS
And all interfaces of the interactive dynamics of PHmD are summarized in Table II as follows.

<table>
<thead>
<tr>
<th>Source Component (Instance)</th>
<th>Sink Component (Instance)</th>
<th>Interface</th>
</tr>
</thead>
<tbody>
<tr>
<td>User Interface (ui)</td>
<td>Sensor (sensor)</td>
<td>interface_ui_sensor</td>
</tr>
<tr>
<td>Actuator (actuator)</td>
<td>User Interface (ui)</td>
<td>interface_actuator_ui</td>
</tr>
<tr>
<td>WirelessBodyEraNetwork (wben)</td>
<td>MobileDevice (md)</td>
<td>interface_wben_md</td>
</tr>
<tr>
<td>CyberInfrastructure (ci)</td>
<td>WirelessBodyEraNetwork (wben)</td>
<td>interface_ci_wben</td>
</tr>
</tbody>
</table>

Let us think about, for definiteness and without loss of generality, that the software architecture model is given by only considering two components as illustrated in Figure 7.

Figure 7. Software architecture for PHmD (partial).

The components are Device and Evaluator. The connectors are eval_dev and dev_eval send physiological signals to Evaluator and response evaluation result to Device, respectively. When Device sends a request message to dev_eval and waits for an evaluation result, dev_eval and eval_dev are assigned as actors, connecting Device and Evaluator. After dev_eval receives a request, it reports an event to Evaluator and waits for a sign of acknowledgment from Evaluator. Evaluator performs the corresponding operation according to forward message. The Cham specification is composed of molecule algebra \( \Sigma_{Cham} \), reaction rules \( T_1, \ldots, T_7 \) and initial solution \( s_0 \). \( \Sigma_{Cham} \) is specified as follows.

\[
Molecule := \text{Com} | \text{Connection} | \text{Data} \\
\text{Com} := \text{Device} : \text{PC} \mid \text{dev_eval} : \text{PC} \mid \text{eval_dev} : \text{PC} \mid \text{Evaluator} : \text{PC} \\
\text{Connection} := i(\text{Role}) \mid o(\text{Role}) \\
\text{Role} := \text{DEcon} \mid \text{EDcon} \mid \text{dev Evalu \mid eval_dev}
\]

where, \text{dev_eval} and \text{eval_dev} are considered as special connectors. \text{Role} is used to denote communication among components specifying message types as different roles. \( i(\text{Role}) \) denotes an import role and \( o(\text{Role}) \) is an export role. Components of \( \text{PHmD} \) in the scenario are \text{Device} and \text{Evaluator}, and connectors are \text{dev_eval} and \text{eval_dev}.

The initial solution and reaction rules are described as follows.

\[
s_0 = \text{Device} \circ \text{DEcon} \circ (\text{dev_eval}), \\
T_1 = i(\text{DEcon}) \circ o(\text{eval_dev}), \\
T_2 = i(\text{DEcon}) \circ o(\text{eval_dev}), \\
T_3 = i(\text{DEcon}) \circ o(\text{eval_dev}), \\
T_4 = o(\text{EDcon}) \circ (\text{dev_eval}), \\
T_5 = o(\text{EDcon}) \circ (\text{eval_dev}), \\
T_6 = o(\text{EDcon}) \circ (\text{eval_dev}), \\
T_7 = o(\text{EDcon}) \circ (\text{eval_dev}), \\
T_8 = o(\text{EDcon}) \circ (\text{eval_dev}), \\
T_9 = o(\text{EDcon}) \circ (\text{eval_dev})
\]

\( T_j \) allows devices to send a request to evaluation center. \( T_2 \) denotes devices are ready to communicate with dev_eval. After sending requests, devices wait for acknowledgment from dev_eval. To generate multiple instances of Device, Cham regenerates Device molecules to allow other Devices to send request messages. The rule restricts some Device sends new requests before receiving acknowledgement from dev_eval. \( T_3 \) specifies the synchronization between devices and dev_eval. After that, dev_eval is ready to receive another device requests. \( T_4 \) specifies eval_dev receives requests from devices and forwards them to evaluate components and after that, eval_dev is ready to receive another device requests. \( T_7 \) denotes device requests are forwarded to evaluator components. Then Evaluator sends an acknowledgement to eval_dev.
T6 denotes \textit{eval-dev} receives the acknowledgement from \textit{Evaluator} and returns it to devices.

T7 denotes devices receive the acknowledgement from \textit{eval-dev} and are ready to send new device requests.

\textbf{B. Quantitative Analysis Results}

Moving beyond the specification of the software architecture and model checking, we have implemented a prototype of PHmD based on Cham. The architecture-layer entities are implemented in Java using Cham specification. Java-RMI interface is employed to store and retrieve necessary architectural information from design phrase and perform a further model check. The self-adaptive effectiveness and performance of Cham framework’s are two important issues requiring evaluation. We use the jogging scenario of PHmD to illustrate the two aspects of the proposed framework. To demonstrate this, an experiment was conducted on a dedicated testbed under different running conditions with regard to different concurrent threads, user requests, wireless bandwidth and servers, respectively. This experiment improved Viterbi-I algorithm on PHmD system that brought about a certain latency for customers. Our experiment results indicate that the proposed method improved system performance while specific loads used in the experiment, as illustrated in Figure 9 compared with Figure 8. Results are shown for system performance without self-adaptation in Figure 8, where the latency experienced by customers never falls on the desired threshold. Figure 9 shows, on the other hand, that the customers experience performance improvement by falling on the optimal value after the period of time is used to execute the optimization algorithm. It is not unexpectedly revealed that the optimization process has an associated latency. However, this is deserved for the PHmD to be self-adaptive according to the optimization algorithm because the latency falls back into the range (depicted in black dash line in Figure 9) afterwards.

\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{fig9.png}
\caption{Overall latency experienced with self-adaptation using the optimization algorithm.}
\end{figure}

\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{fig8.png}
\caption{Latency experienced without self-adaptation.}
\end{figure}

\textbf{VII. CONCLUSIONS AND FUTURE WORKS}

In recent years, a self-adaptation method based on software architecture has become a new research focus. Software architecture provides a whole guideline to create a target system, including specification of architectural elements, constraint of components and the communication contract between architectural elements and software survival environment. The proposed method overcomes the shortcoming of previous research method based on qualitative analysis. CHAM is employed to get specification of software architecture of CPS. By the improved Viterbi-I algorithm, the system promotes performance.

Software architecture-based method showed the effectiveness of performance improvement over the long term. However, the repair intervals of self-adaptive actions are under consideration. The jitter would be possible if the interval is too short. On the other hand, if it is too long, the effectiveness of the self-adaptive actions could not be demonstrated. This remains a wide-open research issue.

\textbf{ACKNOWLEDGMENT}

This work was supported in part by a grant from Zhejiang Provincial Education Department Research Projects (No. Y201119886).

\textbf{REFERENCES}

\begin{enumerate}
\setlength\itemsep{-1pt}
\item Cheng S-W, Garlan D. "Handling Uncertainty in Autonomic Systems," In International Workshop on Living with Uncertainties (IWLU'07), colocated with the 22nd International Conference on Automated Software Engineering (ASE'07), Atlanta, Georgia, USA, pp. 2007.
\end{enumerate}


Hua Wang was born in Zhejiang, China. He received his Ph.D. degree in computer science from Zhejiang University, China, in 2009.

He is currently an associate professor at School of Information and Electronic Engineering, Zhejiang University of Science and Technology, China. His research interests are in Cyber-physical System, Self-adaptive System and Software Architecture.

Dr. Wang is a senior member of International Association of Computer Science and Information Technology (IACSIT).